

DIFFERENTIAL IN INFRASTRUCTURE NETWORKS

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Abstract. Let $G = (V, E)$ be a graph of order n and let $B(D)$ be the set of vertices in $V \setminus D$ that have a neighbor in the vertex set D . The differential of a vertex set D is defined as $\partial(D) = |B(D)| - |D|$ and the maximum value of $\partial(D)$ for any subset D of V is the differential of G . A set D of vertices of a graph G is said to be a dominating set if every vertex in $V \setminus D$ is adjacent to a vertex in D . G is a dominant differential graph if it contains a ∂ -set which is also a dominating set. This paper is devoted to the computation of differential of wheel, cycle and path-related graphs as infrastructure networks. Furthermore, dominant differential wheel, cycle and path-related types of networks are recognized.

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1. INTRODUCTION

In this paper, simple, finite and undirected graphs without loops and multiple edges are considered. Let $G = (V, E)$ be a graph with vertex set V and edge set E . The order of G is given by $|V(G)| = n$ and the size is defined as $|E(G)| = m$ where $|*|$ denotes the number of elements in the set (*i.e.* the cardinality). The neighborhood of a vertex $v \in V(G)$ is the set of vertices adjacent to v , denoted $N_G(v)$ or just $N(v)$, and the closed neighborhood of v is given by $N[v] = N(v) \cup \{v\}$. Thus, $N(v) = \{u \in V(G) | uv \in E(G)\}$ and $N(v)$ is referred to as the open neighborhood of v . The degree of a vertex $v \in V$ is defined as $d(v) = |N(v)|$. For a set $S \subseteq V$, $N(S) = \bigcup_{v \in S} N(v)$ and $N[S] = N(S) \cup S$. An end-vertex or a pendant or pendent vertex is a vertex of degree one and its neighbor is called a support vertex. For $S \subseteq V(G)$, the subgraph of G induced by S is denoted by $G[S]$. Let G and H be two disjoint graphs. The join of graphs G and H , denoted by $G \vee H$, is obtained from the disjoint union G and H by adding the edges $\{xy | x \in V(G), y \in V(H)\}$ [21].

For any real number x we define the ceiling function $\lceil x \rceil$ as the smallest integer greater than or equal to x and similarly we define the floor function $\lfloor x \rfloor$ as the largest integer smallest than or equal to x .

Graph theoretic techniques provide a convenient tool for the investigation of networks. It is well-known that an interconnection network can be modeled by a graph with vertices representing sites of the network and edges representing links between sites of the network. Therefore various network problems can be studied by graph theoretical methods.

The differential in graphs is a subject of increasing interest, both in pure and applied mathematics. The research and application of the $\partial(G)$ appears mainly in computational mathematics. The differential of a graph

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was introduced in [17] in 2006, and studied by several authors [1–9, 14, 16, 19, 20], motivated by its applications to information diffusion in social networks. The study of the mathematical properties of the differential in graphs stated in [1–9, 14, 16, 17, 19, 20]. This parameter has been studied by many authors, both from the viewpoint of combinatorics and from the viewpoint of the algorithmic complexity. We refer to the papers [1–9, 14, 16, 17, 19, 20] and the literature quoted therein. Since computing the differential of a graph is NP-complete in general, it becomes an interesting question to calculate differential for some special classes of interesting or practically useful graphs. In the following sections we will deal with this question.

Let $G = (V, E)$ be a graph of order n , for every set $D \subseteq V$ let $B(D)$ be the set of vertices in $V \setminus D$ that have a neighbor in the vertex set D . The differential of D is defined as $\partial(D) = |B(D)| - |D|$ and the differential of a graph G , written $\partial(G)$, is equal to $\max\{\partial(D) : D \subseteq V\}$. We will say that $D \subseteq V$ is a differential set or ∂ -set if $\partial(D) = \partial(G)$ is called a ∂ -set or differential set. Note that the connectivity of G is not an important restriction, since if G has connected components G_1, \dots, G_k , then $\partial(G) = \partial(G_1) + \dots + \partial(G_k)$. Therefore, we will only consider connected graphs.

A set D of vertices of a graph G is said to be a dominating set if every vertex in $V \setminus D$ is adjacent to a vertex in D . The domination number of G , denoted by $\gamma(G)$ is the minimum size of a dominating set of G [21]. Research on domination in graphs has not only important theoretical signification, but also varied application in such fields as computer science, communication networks, ad hoc networks, biological and social networks, distributed computing, coding theory, and web graphs. Dominating sets in graphs are natural models for facility location problems in operations research. In general, the concept of dominating sets in graph theory finds wide applications in different types of communication networks. A broadcast from a communication vertex is received by all its neighbors. This is captured by the notion of domination in a graph. Finally, we will say that G is a dominant differential graph if it contains a ∂ -set which is also a dominating set. Some examples of dominant differential graphs are complete graphs, star graphs, wheel graphs, and path graphs P_n and cycle graphs C_n with $n = 3k$ or $n = 3k + 2$.

The rest of the paper is structured as follows. In Section 2, the known results in literature are overviewed. In the following sections, the differential of wheel, cycle and path-related types of networks are computed and exact formulae are derived.

2. KNOWN RESULTS

Theorem 2.1 ([9]). *The differential of*

- (a) *the complete graph K_n of order n is $\partial(K_n) = n - 2$;*
- (b) *the path P_n of order n is $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$;*
- (c) *the cycle C_n ($n \geq 3$) of order n is $\partial(C_n) = \lfloor \frac{n}{3} \rfloor$;*
- (d) *the star $K_{1,n}$ of order $n + 1$ is $\partial(K_{1,n}) = n - 1$;*
- (e) *the complete bipartite graph $K_{m,n}$ of order $m + n$ is $\partial(K_{m,n}) = \max\{m - 1, n - 1, m + n - 4\}$;*
- (f) *the wheel W_n of order $n + 1$ is $\partial(W_n) = n - 1$.*

Theorem 2.2 ([7]). *A graph G is dominant differential if and only if $\partial(G) = n - 2\gamma(G)$.*

3. DIFFERENTIAL IN WHEEL RELATED NETWORKS

In this section, the differential of wheel-related networks including gear and helm networks are calculated (Fig. 1).

3.1. Gear networks

Gear network is a wheel graph with a vertex added between each pair adjacent graph vertices of the outer cycle. G_n has $2n + 1$ vertices and $3n$ edges [12]. G_n includes an even cycle C_{2n} . There are two types of vertices

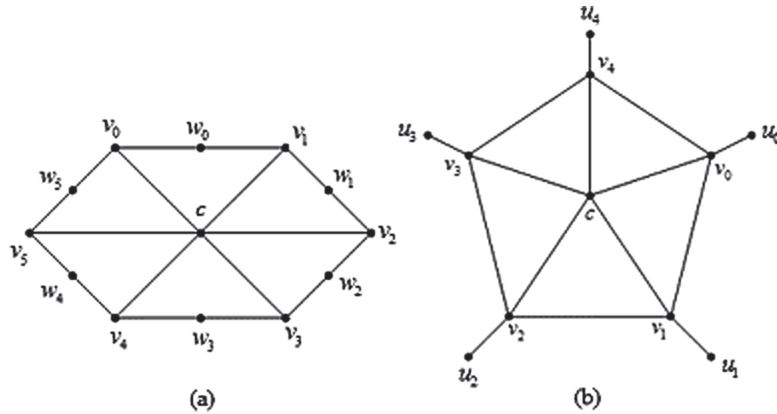


FIGURE 1. (a) Gear network G_n for $n = 6$. (b) Helm network H_n for $n = 5$.

of C_{2n} in G_n as vertices of degree two and three, respectively. The vertices of degree two are referred to as minor vertices and vertices of degree three to as major vertices [15]. The central vertex c of G_n has degree of n . Label the major and minor vertices, respectively, as v_0, \dots, v_{n-1} and w_0, \dots, w_{n-1} and let w_i be adjacent to the vertices v_i and v_{i+1} for $0 \leq i \leq n - 1$, where $i + 1$ is taken modulo n .

Theorem 3.1. *The differential of the gear network G_n of order $2n + 1$ is*

$$\partial(G_n) = \begin{cases} 2 + \lfloor \frac{2n-3}{3} \rfloor, & \text{if } n = 3; \\ n - 1, & \text{if } n \geq 4. \end{cases}$$

Proof. If we take the central vertex c and so $D_1 = \{c\}$, then we have that $B\{D_1\} = \{v_0, \dots, v_{n-1}\}$ and so $\partial(D_1) = n - 1$, and taking any other subset of $V(G_n)$ to the set D_1 yields $\partial(D_1) \leq n - 1$.

If we take a major vertex v_i ($0 \leq i \leq n - 1$) of G_n to the set D_2 , that is $D_2 = \{v_i\}$, then we have $B(D_2) = \{c, w_{i-1}, w_i\}$, where $i - 1$ is taken modulo n , yielding $\partial(D_2) = 2$.

Let $S_1 = V(G_n) \setminus N_{G_n}[v_i]$ and so we have that $G_n[S_1] = C_{2n-3}$. If we take the maximal ∂ -set of $G_n[S_1]$ to the set D_2 having the set D_3 , then since $\partial(C_n) = \lfloor \frac{n}{3} \rfloor$, we receive $\partial(D_3) = \partial(D_1) + \partial(C_{2n-3}) = 2 + \lfloor \frac{2n-3}{3} \rfloor$, and taking any other subset of $V(G_n)$ to the set D_3 yields $\partial(D_3) < 2 + \lfloor \frac{2n-3}{3} \rfloor$.

If we take a minor vertex w_i ($0 \leq i \leq n - 1$) of G_n to the set D_4 , that is $D_4 = \{w_i\}$, then we have $B(D_4) = \{v_i, v_{i+1}\}$, where $i + 1$ is taken modulo n , yielding $\partial(D_4) = 1$.

Let $S_2 = V(G_n) \setminus \{c, N_{G_n}[w_i]\}$ and so we have that $G_n[S_2] = C_{2n-3}$. If we take the maximal ∂ -set of $G_n[S_2]$ to the set D_4 having the set D_5 , then since $\partial(C_n) = \lfloor \frac{n}{3} \rfloor$ and the maximal ∂ -set of $G_n[S_2]$ includes at least one major vertex, we receive $\partial(D_5) = \partial(D_4) + (\partial(C_{2n-3}) + 1) = 2 + \lfloor \frac{2n-3}{3} \rfloor$, and taking any other subset of $V(G_n)$ to the set D_5 yields $\partial(D_5) < 2 + \lfloor \frac{2n-3}{3} \rfloor$.

By the definition of graph differential, among all the differential sets, we get

$$\begin{aligned} \partial(G_n) &= \max \{ \partial(D_k) \} \quad (1 \leq k \leq 5) \\ \partial(G_n) &= \begin{cases} 2 + \lfloor \frac{2n-3}{3} \rfloor, & \text{if } n \leq 6; \\ n - 1, & \text{if } n \geq 4. \end{cases} \end{aligned}$$

Thus, the proof holds. □

Remark 3.2. We can easily observe that $\gamma(G_n) = \lceil \frac{2n}{3} \rceil$ and by Theorem 2.2 we conclude that gear networks are dominant differential for $n = 3, 4, 6$.

3.2. Helm networks

Helm H_n is a network of order $2n+1$ obtained from a wheel W_n with cycle C_n having a pendant edge attached to each vertex of the cycle. H_n consists of the vertex set $V(H_n) = \{v_i | 0 \leq i \leq n-1\} \cup \{u_i | 0 \leq i \leq n-1\} \cup \{c\}$ and edge set $E(H_n) = \{v_i v_{i+1} | 0 \leq i \leq n-1\} \cup \{v_i u_i | 0 \leq i \leq n-1\} \cup \{v_i c | 0 \leq i \leq n-1\}$, where $i+1$ is taken modulo n [12]. The central vertex c of H_n has a vertex degree of n . There are two types of vertices in $H_n \setminus \{c\}$ as the vertices of degree four and one, respectively. The vertices of degree one and four are referred to as minor and major vertices, respectively [15].

Theorem 3.3. *The differential of the helm network H_n of order $2n+1$ is*

$$\partial(H_n) = \begin{cases} 3 + 2 \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3; \\ n - 1, & \text{if } n \geq 4. \end{cases}$$

Proof. If we take the central vertex c of H_n to the set D_1 , then we have $B\{D_1\} = \{v_0, \dots, v_{n-1}\}$ yielding $\partial(D_1) = n - 1$, and taking any other subset of $V(H_n)$ to the set D_1 yields $\partial(D_1) < n - 1$.

If we take a major vertex v_i ($0 \leq i \leq n-1$) of H_n to the set D_2 , that is $D_2 = \{v_i\}$, then we have $B(D_2) = \{c, u_i, v_{i+1}, v_{i-1}\}$, where $i+1$ and $i-1$ are taken modulo n , yielding $\partial(D_2) = 3$.

Let $S_1 = V(H_n) \setminus N_{H_n}[v_i]$ and so we have that $H_n[S_1] = P_{n-3}^*$ where P_n^* is the path graph of order n with a pendant vertex attached to each vertex of the path. If we take the maximal ∂ -set of P_{n-3} to the set D_2 having the set D_3 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$ and every vertex of P_{n-3} is adjacent to a pendant vertex, we receive

$$\partial(D_3) = \begin{cases} \partial(D_2) + \partial(P_{n-3}) + \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3k \text{ or } n = 3k + 1; \\ \partial(D_2) + \partial(P_{n-3}) + \lceil \frac{n-3}{3} \rceil, & \text{if } n = 3k + 2, \end{cases}$$

$$\partial(D_3) = \begin{cases} 3 + 2 \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3k \text{ or } n = 3k + 1; \\ 3 + \lfloor \frac{n-3}{3} \rfloor + \lceil \frac{n-3}{3} \rceil, & \text{if } n = 3k + 2, \end{cases}$$

where $k \in \mathbb{Z}$ and taking any other subset of $V(H_n)$ to the set D_3 , yields

$$\partial(D_3) < \begin{cases} 3 + 2 \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3k \text{ or } n = 3k + 1; \\ 3 + \lfloor \frac{n-3}{3} \rfloor + \lceil \frac{n-3}{3} \rceil, & \text{if } n = 3k + 2. \end{cases}$$

If we take a minor vertex u_i ($0 \leq i \leq n-1$) of H_n to the set D_4 , then we have $B(D_4) = \{v_i\}$, yielding $\partial(D_4) = 0$.

Let $S_2 = V(H_n) \setminus N_{H_n}[u_i]$ and so we have the graph $H_n[S_2]$ including the central vertex c , $n-1$ major and $n-1$ minor vertices of H_n , and also $n-1$ major vertices induce the subgraph P_{n-1} in $H_n[S_2]$. If we take the maximal ∂ -set of P_{n-1} to the set D_4 having the set D_5 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$ and every major vertex is adjacent to a pendant vertex and the central vertex c , we receive

$$\partial(D_5) = \begin{cases} \partial(D_4) + \partial(P_{n-1}) + 1 + \lfloor \frac{n-1}{3} \rfloor, & \text{if } n = 3k + 1 \text{ or } n = 3k + 2; \\ \partial(D_4) + \partial(P_{n-1}) + 1 + \lceil \frac{n-1}{3} \rceil, & \text{if } n = 3k, \end{cases}$$

$$\partial(D_5) = \begin{cases} 2 \lfloor \frac{n-1}{3} \rfloor + 1, & \text{if } n = 3k + 1 \text{ or } n = 3k + 2; \\ \lfloor \frac{n-1}{3} \rfloor + \lceil \frac{n-1}{3} \rceil, & \text{if } n = 3k, \end{cases}$$

where $k \in \mathbb{Z}$ and taking any other subset of $V(H_n)$ to the set D_5 yields

$$\partial(D_5) < \begin{cases} 2 \lfloor \frac{n-1}{3} \rfloor + 1, & \text{if } n = 3k + 1 \text{ or } n = 3k + 2; \\ \lfloor \frac{n-1}{3} \rfloor + \lceil \frac{n-1}{3} \rceil + 1, & \text{if } n = 3k. \end{cases}$$

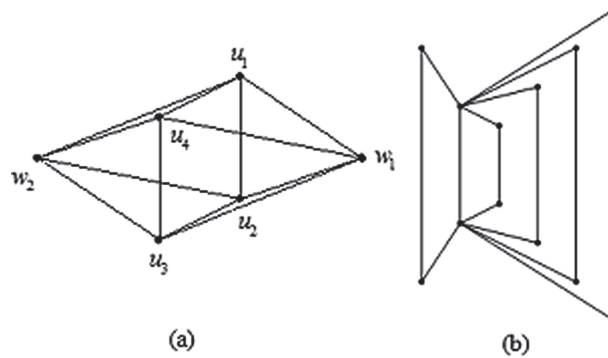


FIGURE 2. (a) Bipyramid network $BP(n)$. (b) n -gon book of k -pages network $B(n, k)$ for $n = 4, n = 5$.

By the definition of graph differential, among all the differential sets, we get

$$\begin{aligned} \partial(H_n) &= \max \{ \partial(D_l) \} \quad (1 \leq l \leq 5) \\ \partial(H_n) &= \begin{cases} 3 + 2 \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3; \\ n - 1, & \text{if } n \geq 4. \end{cases} \end{aligned}$$

Thus, the proof holds. □

4. DIFFERENTIAL IN CYCLE RELATED NETWORKS

In this section, the differential of cycle-related graphs including k -pyramid and n -gon book of k -pages networks are calculated (Fig. 2).

4.1. k -pyramids

The join graph $C_n \vee N_k$ ($n \geq 3, k \geq 1$), where N_k is the null graph of order k , is called k -pyramid and is denoted by $kP(n)$. The 2-pyramid network $C_n \vee N_2$ is called bipyramid network and is denoted by $BP(n)$. The 1-pyramid network $C_n \vee N_1$ is the wheel graph W_n [12].

Let $u_1, u_2, u_3, \dots, u_n$ be the vertices of C_n and $w_1, w_2, w_3, \dots, w_k$ be the vertices of N_k . Then, we have $\deg(u_i) = k + 2$ ($1 \leq i \leq n$) and $\deg(w_j) = n$ ($1 \leq j \leq k$).

Theorem 4.1. *The differential of the k -pyramid network $kP(n)$ of order $n + k$ is*

$$\partial(kP(n)) = \begin{cases} n - 1, & \begin{cases} \text{if } n = 3t & \text{and } n \geq \frac{3k+3}{2}; \\ \text{if } n = 3t + 1 & \text{and } n \geq \frac{3k+2}{2}; \\ \text{if } n = 3t + 2 & \text{and } n \geq \frac{3k+1}{2}; \end{cases} \\ \lfloor \frac{n}{3} \rfloor + k, & \text{otherwise,} \end{cases}$$

where $t \in \mathbb{Z}$.

Proof. If we take a vertex w_j ($1 \leq j \leq k$) of N_k in $kP(n)$ to the set D_1 , that is $D_1 = \{w_j\}$, then we have $B(D_1) = \{u_1, \dots, u_n\}$ and so $\partial(D_1) = n - 1$, and taking any other subset of $V(kP(n))$ to the set D_1 yields $\partial(D_1) < n - 1$.

If we take the maximal ∂ -set of C_n in $kP(n)$ to the set D_2 , then since $\partial(C_n) = \lfloor \frac{n}{3} \rfloor$ and every vertex u_i ($1 \leq i \leq n$) of C_n in $kP(n)$ is adjacent to every vertex w_j ($1 \leq j \leq k$) of N_k , we have $\partial(D_2) = \lfloor \frac{n}{3} \rfloor + k$, and taking any other subset of $V(kP(n))$ to the set D_2 yields $\partial(D_2) < \lfloor \frac{n}{3} \rfloor + k$.

By the definition of graph differential, we have

$$\begin{aligned} \partial(kP(n)) &= \max \{ \partial(D_p) \} \quad (p = 1, 2) \\ \partial(kP(n)) &= \begin{cases} n - 1, & \begin{cases} \text{if } n = 3t & \text{and } n \geq \frac{3k+3}{2}; \\ \text{if } n = 3t + 1 & \text{and } n \geq \frac{3k+2}{2}; \\ \text{if } n = 3t + 2 & \text{and } n \geq \frac{3k+1}{2}; \end{cases} \\ \lfloor \frac{n}{3} \rfloor + k, & \text{otherwise,} \end{cases} \end{aligned}$$

where $t \in \mathbb{Z}$. Thus, the proof holds. □

Corollary 4.2. *The differential of the bipyramid network $BP(n)$ of order $n + 2$ is*

$$\partial(BP(n)) = \begin{cases} \lfloor \frac{n}{3} \rfloor + 2, & \text{if } n = 3; \\ n - 1, & \text{if } n > 3. \end{cases}$$

Remark 4.3. Since the wheel network W_n is 1-pyramid network $C_n \vee N_1$, by taking $k = 1$, the value of $\partial(W_n)$ in Theorem 2.1(f) holds.

Remark 4.4. We can easily observe that $\gamma(kP(n)) = k$ and by Theorem 2.2 we conclude that k -pyramids are dominant differential networks for

$$\begin{cases} n = 3t & \text{and } n = 3k; \\ n = 3t + 1 & \text{and } n = \frac{6k-1}{2}; \\ n = 3t + 2 & \text{and } n = 3k - 1, \end{cases}$$

where $t \in \mathbb{Z}$.

4.2. n -gon books

When k copies of C_n ($n \geq 3$) share a common edge, it will form an n -gon book of k pages and is denoted by $B(n, k)$. The degree set of $B(n, k)$ is $\{2, k + 1\}$ [12]. Therefore, the vertices of $B(n, k)$ are of two kinds: vertices of degree 2 are referred to as minor vertices and vertices of degree $k + 1$ to as major vertices. For $k = 1$, we notice that $B(n, k) \cong C_n$.

Theorem 4.5. *The differential of the n -gon book of k ($k > 1$) pages network $B(n, k)$ of order $(n - 2)k + 2$ is*

$$\partial(B(n, k)) = \begin{cases} \frac{nk}{3}, & \text{if } n = 3t; \\ k \left(\lfloor \frac{n-4}{3} \rfloor + 2 \right) - 2, & \text{if } n \neq 3t, \end{cases}$$

where $t \in \mathbb{Z}$.

Proof. If we take one of the major vertices of $B(n, k)$ -say vertex u , to the set D_1 , that is $D_1 = \{u\}$, then since $|N_{B(n,k)}(u)| = k + 1$, we have $\partial(D_1) = k$.

Let $S_1 = V(B(n, k)) \setminus N_{B(n,k)}[u]$ and so we have that $B(n, k)[S_1] = \bigcup_{i=1}^k P_{n-3}$. If we take the maximal ∂ -set of $B(n, k)[S_1]$ to the set D_1 having the set D_2 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_2) = k \lfloor \frac{n-3}{3} \rfloor + k$.

For $n = 3t + 1$ ($t \in \mathbb{Z}$), if we take the other major vertex of $B(n, k)$ -say vertex v to the set D_2 having the set D_3 , then we have $\partial(D_3) \leq (k \lfloor \frac{n-3}{3} \rfloor + k) + (k - 2) = k \lfloor \frac{n-3}{3} \rfloor + 2k - 2$, and taking any other subset of $V(B(n, k))$ to the set D_3 yields $\partial(D_3) < k \lfloor \frac{n-3}{3} \rfloor + 2k - 2$.

If we take $D_4 = \{u, v\}$ as the set of major vertices, then we have $\partial(D_4) = 2k - 2$.

For $n \neq 3t + 1$ ($t \in \mathbb{Z}$), let $S_2 = V(B(n, k)) \setminus N_{B(n,k)}[D_3]$ and so we have that $B(n, k)[S_2] = \bigcup_{i=1}^k P_{n-4}$. If we take the maximal ∂ -set of $B(n, k)[S_2]$ to the set D_4 having the set D_5 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_5) = 2k - 2 + k \lfloor \frac{n-4}{3} \rfloor$, and taking any other subset of $V(B(n, k))$ yields $\partial(D_5) < 2k - 2 + k \lfloor \frac{n-4}{3} \rfloor$.

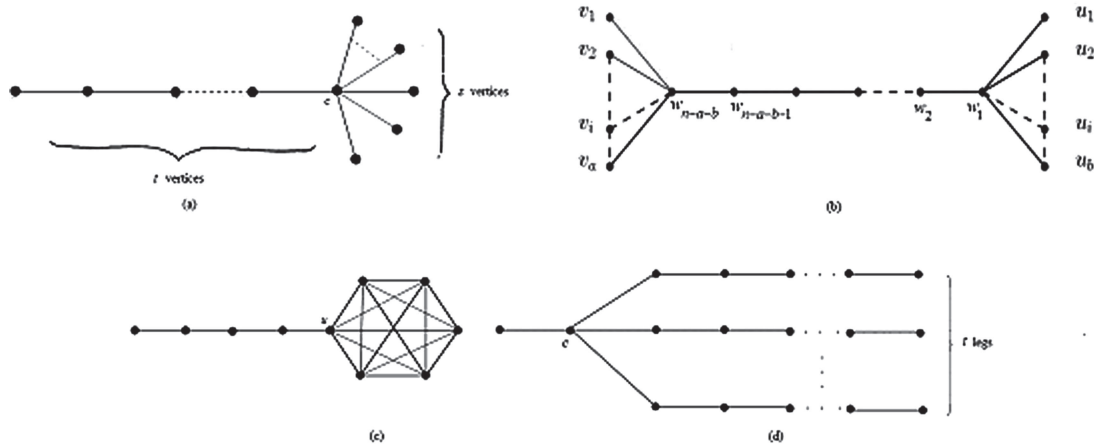


FIGURE 3. (a) Comet network $C_{s,t}$. (b) Double comet network $DC(n, a, b)$. (c) Lollipop networks $L_{n,d}$ for $n = 10$ and $d = 5$. (d) E_p^t network.

By the definition of graph differential, among the differential sets, we receive for $k > 1$

$$\partial(B(n, k)) = \max \{ \partial(D_i) \} \quad (1 \leq i \leq 5)$$

$$\partial(B(n, k)) = \begin{cases} k \left(\lfloor \frac{n-3}{3} \rfloor + 1 \right), & \text{if } n = 3t; \\ k \left(\lfloor \frac{n-4}{3} \rfloor + 2 \right) - 2, & \text{if } n \neq 3t, \end{cases}$$

where $t \in \mathbb{Z}$. The theorem is thus proved. □

Remark 4.6. We can easily observe that

$$\gamma(B(n, k)) = \begin{cases} 1 + k \lfloor \frac{n-3}{3} \rfloor, & \text{if } n = 3t \text{ or } n = 3t + 2; \\ 2 + k \lfloor \frac{n-4}{3} \rfloor, & \text{if } n = 3t + 1, \end{cases}$$

where $t \in \mathbb{Z}$ and by Theorem 2.2 we conclude that n -gon books are dominant differential networks for $n = 3t + 2$ and $k = 2$ or $n = 3t$ ($t \in \mathbb{Z}$).

5. DIFFERENTIAL IN PATH RELATED NETWORKS

In this section, the differential of path-related networks including comet, double comet, lollipop and E_p^t networks are calculated (Fig. 3).

5.1. Comet networks

The comet $C_{s,t}$ where s and t are positive integers, denotes the tree obtained by identifying the center of the star $K_{1,s}$ with an end-vertex of P_t , the path of order t . So $C_{s,1} \cong K_{1,s}$ and $C_{1,p-1} \cong P_p$ [13]. Let the center of the star $K_{1,s}$ - that is one end-vertex of P_t be the vertex c .

Theorem 5.1. *The differential of the comet network $C_{s,t}$ of order $s + t$ is*

$$\partial(C_{s,t}) = s + \left\lfloor \frac{t-2}{3} \right\rfloor.$$

Proof. If we take the center vertex c of $K_{1,s}$ in $C_{s,t}$ to the set D_1 , that is $D_1 = \{c\}$, then we receive $\partial(D_1) = 5$.

Let $S_1 = V(C_{s,t}) \setminus N_{C_{s,t}}[c]$ and so we have that $C_{s,t}[S_1] = P_{t-2}$. If we take the maximal ∂ -set of $C_{s,t}[S_1]$ to the set D_1 having the set D_2 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_2) = s + \lfloor \frac{t-2}{3} \rfloor$, and taking any other subset of $V(C_{s,t})$ to the set D_2 yields $\partial(D_2) < s + \lfloor \frac{t-2}{3} \rfloor$.

If we take one of the vertices of $K_{1,s}$ except the center vertex c in $C_{s,t}$ to the set D_3 , then we have $\partial(D_3) = 0$.

Let $S_2 = V(C_{s,t}) \setminus N_{C_{s,t}}[D_3]$ and so we have that $C_{s,t}[S_2] = P_{t-1}$. If we take the maximal ∂ -set of $C_{s,t}[S_2]$ to the set D_3 having the set D_4 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_4) = \lfloor \frac{t-1}{3} \rfloor$, and taking any other subset of $V(C_{s,t})$ to the set D_4 yields $\partial(D_4) < \lfloor \frac{t-1}{3} \rfloor$.

If we take the maximal ∂ -set of P_t to the set D_5 such that $c \in B(D_5)$, then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_5) = \lfloor \frac{t}{3} \rfloor$.

If we take the center vertex c of $K_{1,s}$ in $C_{s,t}$ to the set D_5 having the set D_6 , then we have $\partial(D_6) = \lfloor \frac{t}{3} \rfloor + s - 2$, and taking any other subset of $V(C_{s,t})$ to the set D_6 yields $\partial(D_6) < \lfloor \frac{t}{3} \rfloor + s - 2$.

By the definition of graph differential, we get

$$\begin{aligned} \partial(C_{s,t}) &= \max \{ \partial(D_i) \} \quad (1 \leq i \leq 6) \\ \partial(C_{s,t}) &= s + \left\lfloor \frac{t-2}{3} \right\rfloor. \end{aligned}$$

Thus the proof holds. □

Remark 5.2. We can easily observe that $\gamma(C_{s,t}) = 1 + \lceil \frac{t-2}{3} \rceil$ and by Theorem 2.2 we conclude that comets are dominant differential networks for $t = 3k + 2$ or $t = 3k + 4$ ($k \in \mathbb{Z}$).

5.2. Double comet networks

For $a, b > 1, n \geq a + b + 2$ by $DC(n, a, b)$ we denote a double comet, which is a tree composed of a path containing $n - a - b$ vertices with a pendent vertices attached to one of the ends of the path and b pendent vertices attached to the other end of the path. Thus, $DC(n, a, b)$ has n vertices and $a + b$ leaves [11]. Let $v_1, v_2, \dots, v_a, u_1, u_2, \dots, u_b, w_1, w_2, \dots, w_{n-a-b}$ be the vertex set of the double comet $DC(n, a, b)$, which is obtained from a path P_{n-a-b} of vertices $w_1, w_2, \dots, w_{n-a-b}$ by attaching the pendant vertices u_1, u_2, \dots, u_b to the one end vertex w_1 of P_{n-a-b} and attaching the pendant vertices v_1, v_2, \dots, v_a to the other end vertex w_n of P_{n-a-b} .

Theorem 5.3. *The differential of the double comet network $DC(n, a, b)$ of order n ($a, b > 1, n - a - b > 2$) is*

$$\partial(DC(n, a, b)) = a + b + \left\lfloor \frac{n - a - b - 4}{3} \right\rfloor.$$

Proof. If we take the vertex w_1 of $DC(n, a, b)$ to the set D_1 , that is $\partial(D_1) = a$.

Let $S_1 = V(DC(n, a, b)) \setminus N_{DC(n, a, b)}[w_1]$ and so we have that $DC(n, a, b)[S_1] = C_{b, n-a-b-2}$, where $C_{b, n-a-b-2}$ is a comet graph of order $n - a - 2$. If we take the maximal ∂ -set of $DC(n, a, b)[S_1]$ to the set D_1 having the set D_2 , then since $\partial(C_{s,t}) = s + \lfloor \frac{t-2}{3} \rfloor$, we receive $\partial(D_2) = a + b + \lfloor \frac{n-a-b-4}{3} \rfloor$, and taking any other subset of $V(DC(n, a, b))$ to the set D_2 yields $\partial(D_2) < a + b + \lfloor \frac{n-a-b-4}{3} \rfloor$.

If we take the vertex w_{n-a-b} of $DC(n, a, b)$ to the set D_3 , then we have $\partial(D_3) = b$.

Let $S_2 = V(DC(n, a, b)) \setminus N_{DC(n, a, b)}[w_n]$ and so we have that $DC(n, a, b)[S_2] = C_{a, n-a-b-2}$. If we take the maximal ∂ -set of $DC(n, a, b)[S_2]$ to the set D_3 having the set D_4 , then since $\partial(C_{s,t}) = s + \lfloor \frac{t-2}{3} \rfloor$, we have $\partial(D_4) = b + a + \lfloor \frac{n-a-b-4}{3} \rfloor$, and taking any other subset of $V(DC(n, a, b))$ to the set D_4 yields $\partial(D_4) < b + a + \lfloor \frac{n-a-b-4}{3} \rfloor$.

If we take the vertex v_i ($1 \leq i \leq a$) of $DC(n, a, b)$ to the set D_5 , then we have $\partial(D_5) = 0$.

Let $S_3 = V(\text{DC}(n, a, b)) \setminus N_{\text{DC}(n, a, b)}[v_i]$ and so we have that $\text{DC}(n, a, b)[S_3] = \overline{K_{a-1}} \cup C_{b, n-a-b-1}$. If we take the maximal ∂ -set of $\text{DC}(n, a, b)[S_3]$ to the set D_5 having the set D_6 , then since $\partial(C_{s,t}) = s + \lfloor \frac{t-2}{3} \rfloor$, we have $\partial(D_6) = \partial(C_{b, n-a-b-1}) = b + \lfloor \frac{n-a-b-3}{3} \rfloor$, and taking any other subset of $V(\text{DC}(n, a, b))$ to the set D_6 yields $\partial(D_6) < b + \lfloor \frac{n-a-b-3}{3} \rfloor$.

If we take the vertex u_j ($1 \leq j \leq b$) of $\text{DC}(n, a, b)$ to the set D_7 , then we have $\partial(D_7) = 0$.

Let $S_4 = V(\text{DC}(n, a, b)) \setminus N_{\text{DC}(n, a, b)}[u_j]$ and so we have that $\text{DC}(n, a, b)[S_4] = \overline{K_{b-1}} \cup C_{a, n-a-b-1}$. If we take the maximal ∂ -set of $\text{DC}(n, a, b)[S_4]$ to the set D_7 having the set D_8 , then since $\partial(C_{s,t}) = s + \lfloor \frac{t-2}{3} \rfloor$, we have $\partial(D_8) = \partial(C_{a, n-a-b-1}) = a + \lfloor \frac{n-a-b-3}{3} \rfloor$, and taking any other subset of $V(\text{DC}(n, a, b))$ to the set D_8 yields $\partial(D_8) < a + \lfloor \frac{n-a-b-3}{3} \rfloor$.

For $n - a - b = 3k + 1$ ($k \in \mathbb{Z}$), if we take the maximal ∂ -set of P_{n-a-b} to the set D_9 such that $w_1 \in B(D_9)$ or $w_{n-a-b} \in B(D_9)$, then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_9) = \lfloor \frac{n-a-b}{3} \rfloor$. If $w_1 \in B(D_9)$, then let $S_5 = V(\text{DC}(n, a, b)) \setminus N_{\text{DC}(n, a, b)}[D_9 \cup B(D_9)]$ and so we have $\text{DC}(n, a, b)[S_5] = \overline{K_b} \cup K_{1,a}$. If we take the maximal ∂ -set of $\text{DC}(n, a, b)[S_5]$ to the set D_9 having the set D_{10} , then since $\partial(K_{1,n}) = n - 1$, we have $\partial(D_{10}) = \lfloor \frac{n-a-b}{3} \rfloor + (a - 1)$. If we take the vertex w_1 of $\text{DC}(n, a, b)$ to the set D_{10} having the set D_{11} , we have $\partial(D_{11}) = \lfloor \frac{n-a-b}{3} \rfloor + (a + b - 3)$, and taking any other subset of $V(\text{DC}(n, a, b))$ to the set D_{11} yields $\partial(D_{11}) < \lfloor \frac{n-a-b}{3} \rfloor + (a + b - 3)$. If $w_{n-a-b} \in B(D_9)$, then let $S_6 = V(\text{DC}(n, a, b)) \setminus N_{\text{DC}(n, a, b)}[D_9 \cup B(D_9)]$ and so we have $\text{DC}(n, a, b)[S_6] = \overline{K_a} \cup K_{1,b}$. If we take the maximal ∂ -set of $\text{DC}(n, a, b)[S_6]$ to the set D_9 having the set D_{10} , then since $\partial(K_{1,n}) = n - 1$, we have $\partial(D_{10}) = \lfloor \frac{n-a-b}{3} \rfloor + (b - 1)$. If we take the vertex w_1 of $\text{DC}(n, a, b)$ to the set D_{10} having the set D_{11} , we have $\partial(D_{11}) = \lfloor \frac{n-a-b}{3} \rfloor + (b + a - 3)$, and taking any other subset of $V(\text{DC}(n, a, b))$ to the set D_{11} yields $\partial(D_{11}) < \lfloor \frac{n-a-b}{3} \rfloor + (b + a - 3)$.

For $n - a - b = 3k$ or $n - a - b = 3k + 2$ ($k \in \mathbb{Z}$), if we take the maximal ∂ -set of P_{n-a-b} to the set D_{12} such that $w_1, w_{n-a-b} \in B(D_{12})$, then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_{12}) = \lfloor \frac{n-a-b}{3} \rfloor$. If we take the vertex $w_1 \in V(\text{DC}(n, a, b))$ to the set D_{12} having the set D_{13} , we have $\partial(D_{13}) = \lfloor \frac{n-a-b}{3} \rfloor + (b - 2)$. If we take the vertex $w_{n-a-b} \in V(\text{DC}(n, a, b))$ to the set D_{13} having the set D_{14} , we have $\partial(D_{14}) = \lfloor \frac{n-a-b}{3} \rfloor + (b + a - 4)$, and taking any subset of $V(\text{DC}(n, a, b))$ to the set D_{14} yields $\partial(D_{14}) < \lfloor \frac{n-a-b}{3} \rfloor + (b + a - 4)$. If we take the vertex $w_{n-a-b} \in V(\text{DC}(n, a, b))$ to the set D_{12} having the set D_{15} , then we have $\partial(D_{15}) = \lfloor \frac{n-a-b}{3} \rfloor + (a - 2)$.

By the definition of graph differential, we get

$$\begin{aligned} \partial(\text{DC}(n, a, b)) &= \max \{ \partial(D_l) \} \quad (1 \leq l \leq 15) \\ \partial(\text{DC}(n, a, b)) &= a + b + \left\lfloor \frac{n - a - b - 4}{3} \right\rfloor \quad \text{for } a, b > 1 \text{ and } n - a - b > 2. \end{aligned}$$

Thus the proof holds. □

Remark 5.4. We can easily observe that $\gamma(\text{DC}(n, a, b)) = 2 + \lceil \frac{n-a-b-4}{3} \rceil$ for $a, b > 1$ and by Theorem 2.2 we conclude that double comets are dominant differential networks for $n - a - b = 3k$ or $n - a - b = 3k + 1$ ($k \in \mathbb{Z}$).

5.3. Lollipop networks

The lollipop network $L_{n,d}$ is a graph obtained from a complete graph K_{n-d} and a path P_d , by joining one of the end vertices of P_d [18], let this vertex be the vertex u , to all the vertices of K_{n-d} .

Theorem 5.5. *The differential of the lollipop network $L_{n,d}$ ($d > 1$) of order n is*

$$\partial(L_{n,d}) = n - d + \left\lfloor \frac{d - 2}{3} \right\rfloor.$$

Proof. If we take the vertex u of $L_{n,d}$ to the set D_1 , that is $D_1 = \{u\}$, then we have $\partial(D_1) = n - d$.

Let $S_1 = V(L_{n,d}) \setminus N_{L_{n,d}}[u]$ and so we have that $L_{n,d}[S_1] = P_{d-2}$. If we take the maximal ∂ -set of $L_{n,d}[S_1]$ to the set D_1 having the set D_2 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_2) = n - d + \lfloor \frac{d-2}{3} \rfloor$, and taking any other subset of $V(L_{n,d})$ to the set D_2 yields $\partial(D_2) < n - d + \lfloor \frac{d-2}{3} \rfloor$.

If we take one of the vertices of K_{n-d} to the set D_3 , then we have $\partial(D_3) = n - d - 1$.

Let $S_2 = V(L_{n,d}) \setminus N_{L_{n,d}}[D_3]$ and so we have that $L_{n,d}[S_2] = P_{d-1}$. If we take the maximal ∂ -set of $L_{n,d}[S_2]$ to the set D_3 having the set D_4 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_4) = n - d - 1 + \lfloor \frac{d-1}{3} \rfloor$, and taking any other subset of $V(L_{n,d})$ to the set D_4 yields $\partial(D_4) < n - d - 1 + \lfloor \frac{d-1}{3} \rfloor$.

If we take the maximal ∂ -set of P_d to the set D_5 , such that $u \in B(D_5)$, then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_5) = \lfloor \frac{d}{3} \rfloor$.

If we take the ∂ -set of K_{n-d} to the set D_5 having the set D_6 , then since $\partial(K_n) = n - 2$, we get $\partial(D_6) = \lfloor \frac{d}{3} \rfloor + n - d - 2$ and taking any other subset of $V(L_{n,d})$ to the set D_6 yields $\partial(D_6) < \lfloor \frac{d}{3} \rfloor + n - d - 2$.

By the definition of graph differential, we get

$$\begin{aligned} \partial(L_{n,d}) &= \max \{ \partial(D_i) \} \quad (1 \leq i \leq 6) \\ \partial(L_{n,d}) &= n - d + \left\lfloor \frac{d-2}{3} \right\rfloor \quad \text{for } d > 1. \end{aligned}$$

Thus the proof holds. □

Remark 5.6. We can easily observe that $\gamma(L_{n,d}) = 1 + \lceil \frac{d-2}{3} \rceil$ for $d > 1$ and by Theorem 2.2 we conclude that lollipop networks are dominant differential for $d = 3k + 2$ or $d = 3k + 4$ ($k \in \mathbb{Z}$).

5.4. E_p^t networks

The network E_p^t is a tree of order $pt + 2$ obtained from a path with two vertices having one of the end-vertices attached to t legs and each leg has p vertices [10]. Let the end-vertex attached to t legs be the vertex c , and the vertex degree of the vertex c is $\deg(c) = t + 1$.

Theorem 5.7. *The differential of the E_p^t network of order $pt + 2$ is*

$$\partial(E_p^t) = t \left(\left\lfloor \frac{p-1}{3} \right\rfloor + 1 \right).$$

Proof. If we take the vertex c of E_p^t to the set D_1 , then we have $\partial(D_1) = t$.

Let $S_1 = V(E_p^t) \setminus N_{E_p^t}[c]$ and so we have that $E_p^t[S_1] = \bigcup_{i=1}^t P_{p-1}$. If we take the maximal ∂ -set of $E_p^t[S_1]$ to the set D_1 having the set D_2 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_2) = t + t \lfloor \frac{p-1}{3} \rfloor$, and taking any other subset of $V(E_p^t)$ to the set D_2 yields $\partial(D_2) < t + t \lfloor \frac{p-1}{3} \rfloor$.

If we take the other end-vertex that is adjacent to the vertex c in P_2 to the set D_3 , then we have $\partial(D_3) = 0$.

Let $S_2 = V(E_p^t) \setminus N_{E_p^t}[D_3]$ and so we have that $E_p^t[S_2] = \bigcup_{i=1}^t P_p$. If we take the maximal ∂ -set of $E_p^t[S_2]$ to the set D_3 having the set D_4 , then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we receive $\partial(D_4) = t \lfloor \frac{p}{3} \rfloor$, and taking any other subset of $V(E_p^t)$ to the set D_4 yields $\partial(D_4) < t \lfloor \frac{p}{3} \rfloor$.

If we take the maximal ∂ -set of one of the legs to the set D_5 such that $c \in B(D_5)$, then since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we have $\partial(D_5) = \lfloor \frac{p+1}{3} \rfloor$.

If we take the maximal ∂ -sets of $t - 1$ legs of E_p^t to the set D_5 having the set D_6 , since $\partial(P_n) = \lfloor \frac{n}{3} \rfloor$, we receive $\partial(D_6) = \lfloor \frac{p+1}{3} \rfloor + (t - 1) \lfloor \frac{p}{3} \rfloor$, and taking any other subset of $V(E_p^t)$ to the set D_6 yields $\partial(D_6) < t + t \lfloor \frac{p+1}{3} \rfloor$.

By the definition of graph differential, we get

$$\begin{aligned} \partial(E_p^t) &= \max \{ \partial(D_i) \} \quad (1 \leq i \leq 6) \\ \partial(E_p^t) &= t \left(\left\lfloor \frac{p-1}{3} \right\rfloor + 1 \right). \end{aligned}$$

Thus the proof holds. □

Remark 5.8. We can easily observe that $\gamma(E_p^t) = 1 + t \lceil \frac{p-1}{3} \rceil$ and by Theorem 2.2 we conclude that E_p^t networks are dominant differential for $p = 3k$ or $p = 3k + 1$ ($k \in \mathbb{Z}$).

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